Vef MA, Moti N, S["]uß T et al. GekkoFS — A temporary burst buffer file system for HPC applications. JOURNAL OF COMPUTER SCIENCE AND TECHNOLOGY 35(1): 72–91 Jan. 2020. DOI 10.1007/s11390-020-9797-6

GekkoFS

A temporary burst buffer file system for HPC applications

<u>Marc-André Vef*</u>, Nafiseh Moti, Tim Süß*, Markus Tacke, André Brinkmann*, Tommaso Tocci⁺, Alberto Miranda⁺, Ramon Nou⁺, Toni Cortes⁺ §

Barcelona

Center

Supercomputing

Centro Nacional de Supercomputación

- * Johannes Gutenberg University Mainz, Center for Data Processing, Mainz, Germany
- + Barcelona Supercomputing Center, Barcelona, Spain
- § Universitat Politenica de Catalunya, Computer Architecture Department, Barcelona, Spain





DFG Deutsche Forschungsgemeinschaft German Research Foundation





UPC

Departament d'Arquitectura de Computadors

UNIVERSITAT POLITÈCNICA DE CATALUNYA



- HPC is increasingly more used by data-driven science applications
- Data-driven workloads impose new requirements on HPC file systems
 - Many metadata operations, random access, small I/O requests ...
- Solutions are software (e.g., ADIOS) or hardware (burst buffers) based
- Burst buffer file systems use node-local storage to accelerate I/O
 - But, they are often full or near POSIX-compliant

Goal

- Goal: Deploy a lightweight file system per job across all allocated nodes
 - Temporary lifetime, e.g., for HPC job or campaigns
 - Use unused node-local storage (RAMDisk, SSD, NVRAM, ...)
 - Inputs are staged into the FS before job starts (Output vice versa)
 - Relax POSIX semantics, e.g., no sequentialized creates
 - Offer only FS features which are actually required by most (not all) applications
- Application assumptions:
 - Each FS object is accessed by a single application
 - Working data set fits into available node-local storage



File system design

Key properties

- Data and metadata are distributed evenly across all job nodes
 - Key-Value store (one per node) handles metadata
 - Node-local file system is used for data
- Strong consistency for direct operations on FS objects
- Synchronous and no cache mechanisms
- No internode-locking and no permission handling
- File I/O is split into equally sized chunks (configurable)
- The node destination of each chunk is computed on the fly
- Chunks are mapped to files in the node-local file system



Evaluation

- GekkoFS and BeeGFS weakly scaled (100K files per process)
 - More than 819 million files in total with 512 nodes



Number of nodes (16 processes / node)



Evaluation

Sequential I/O throughput vs. BeeGFS



